

**IN THE SPECIFICATION:**

Please amend the specification paragraph on page 2, lines 1 19 as follows:

—There currently exist schemes for reducing operating system process interruptions. For instance, the reference entitled “Design and Analysis of Internal Organizations for Compressed Random Access Memories” by Peter A. Franaszek and John T. Robinson, IBM Research Report RC21146(94535), dated October 28, 1998, describes a low level main memory design for storing compressed data that includes a directory portion and a collection of fixed size blocks which are used to store lines in compressed format. In the memory storage scheme described ~~herein~~ therein, highly compressible lines may be stored entirely within a directory entry; otherwise, the directory entry points to one or more of the fixed size blocks which are used to store the line in compressed format. The system further makes use of page tables which translate virtual addresses to real addresses which correspond to the location in the directory of the directory entry for the line and which includes information pertaining to blocks holding a compressed line. Specifically, the information in a directory entry includes flags, fragment combining information, and, assuming fixed size entry structure, pointers to one or more fixed size blocks. On a cache miss, the memory controller and decompression hardware finds the blocks allocated to store the compressed line and dynamically decompresses the line to handle the miss. Similarly, when a new or modified line is stored, the blocks currently allocated to the line are made free (if the line currently resides in the RAM), the line is compressed, and then stored in the RAM by allocating the required number of blocks. --